

It Wasn't Me, It Was the Prototype!

Towards a Formal Model of JavaScript Prototype Pollution

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KTH



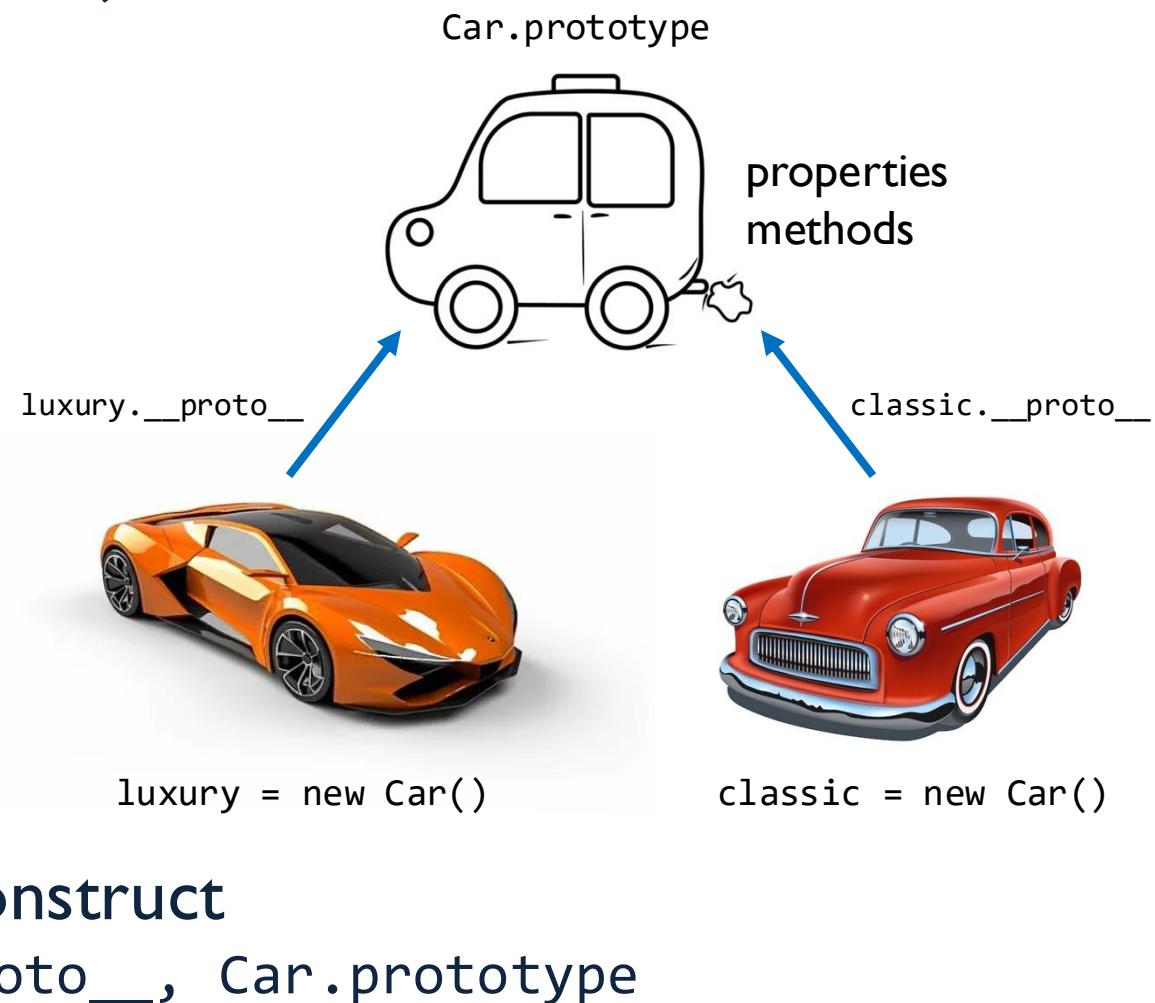
Chalmers

ShiftLeft Workshop, Gothenburg

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JavaScript objects

- **Mutable** collection of properties
 - Values evaluated at **runtime**
- **Prototype: object blueprint**
 - Reusing existing objects
 - An object with a set of properties and functions **shared between all objects of the same type**
 - Exposed as regular programming construct
`luxury["__proto__"], classic.__proto__, Car.prototype`



Prototype-based Inheritance



snowy : Dog	
__proto__	•
age	5

Dog.prototype	
__proto__	•
numOfLegs	4

Dog	
prototype	•

Animal.prototype	
__proto__	•
eat()	function {...}

Animal	
prototype	•

`snowy.age`
`snowy.numOfLegs`
`snowy.eat()`
`snowy.toString()`
`snowy.prop`

Object.prototype	
__proto__	null
toString()	function {...}

Object	
prototype	•



Prototype pollution (PP)

brutalicious : Dog	
__proto__	●
age	100

`brutalicious.__proto__.numOfLegs = 3`

snowy : Dog	
__proto__	●
5	

Poor
snowy! 😞



Dog.prototype	
__proto__	●
numOfLegs	3

Animal.prototype	
__proto__	●
eat()	function {...}

Object.prototype	
__proto__	null
toString()	function {...}

Object	
prototype	●

Dog	
prototype	●

Animal	
prototype	●

Prototype pollution (cont.)

- A vulnerability where an attacker can modify an **object's prototype** at ***runtime***
 - May then be inherited by user-defined objects
 - May then result in **unintended behavior**

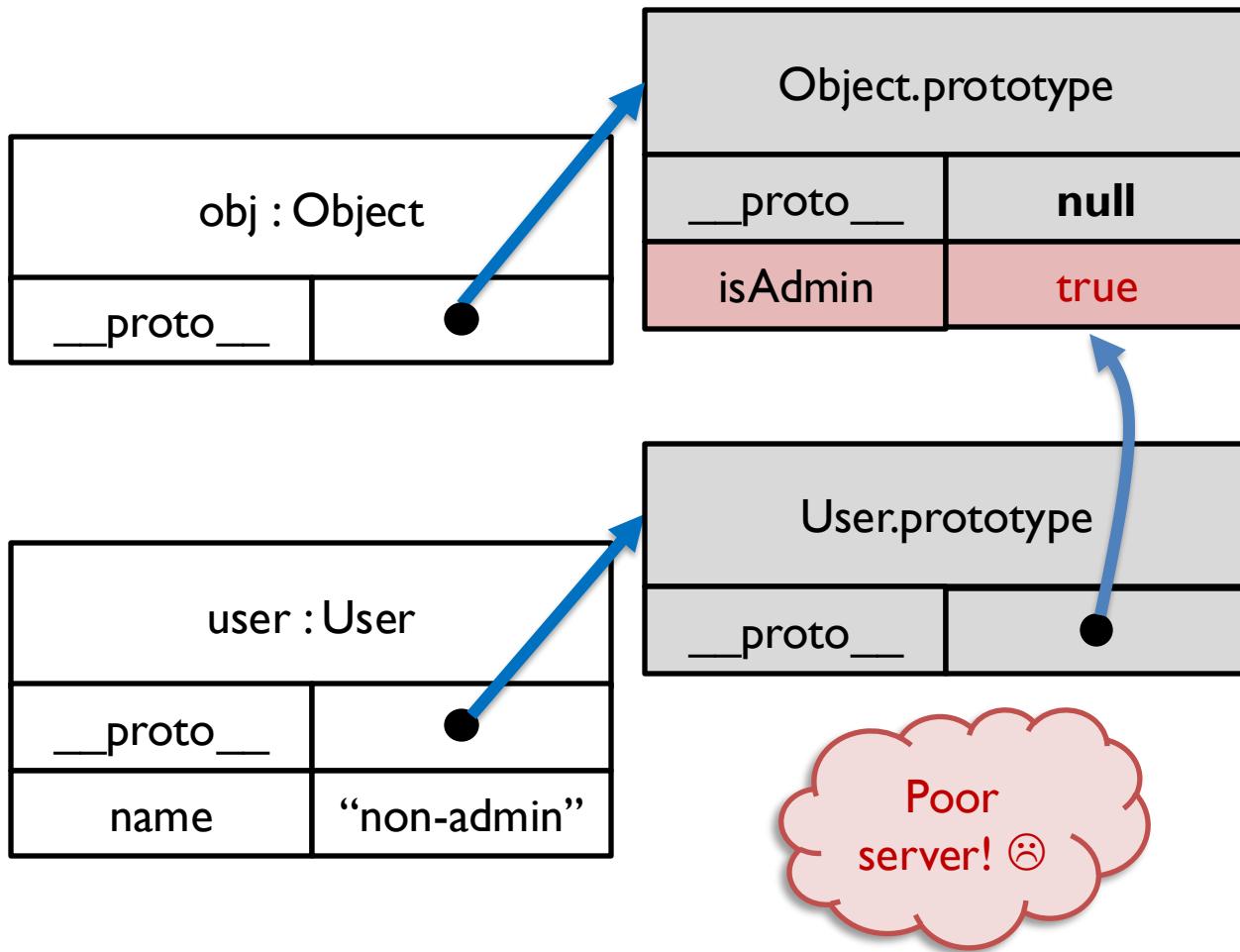


Prototype pollution (cont.)

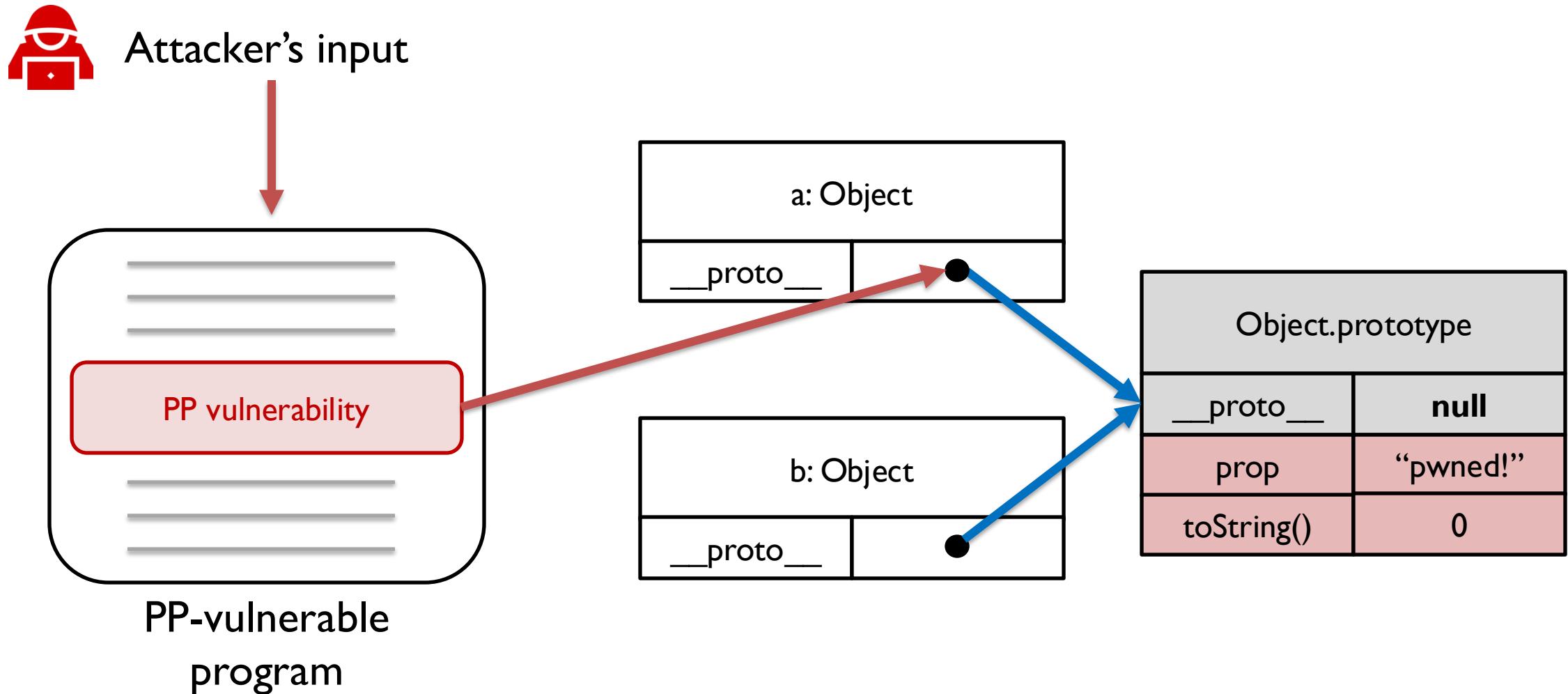
```
{ "__proto__": { "isAdmin": true } }
```



```
function checkAdmin(user) {  
  if (user.isAdmin) {  
    console.log("Welcome, admin user!!");  
  } else {  
    console.log("Just a regular user...");  
  }  
}  
...  
let user = new User("non-admin");  
...  
let obj = JSON.parse(input);  
checkAdmin(user);
```

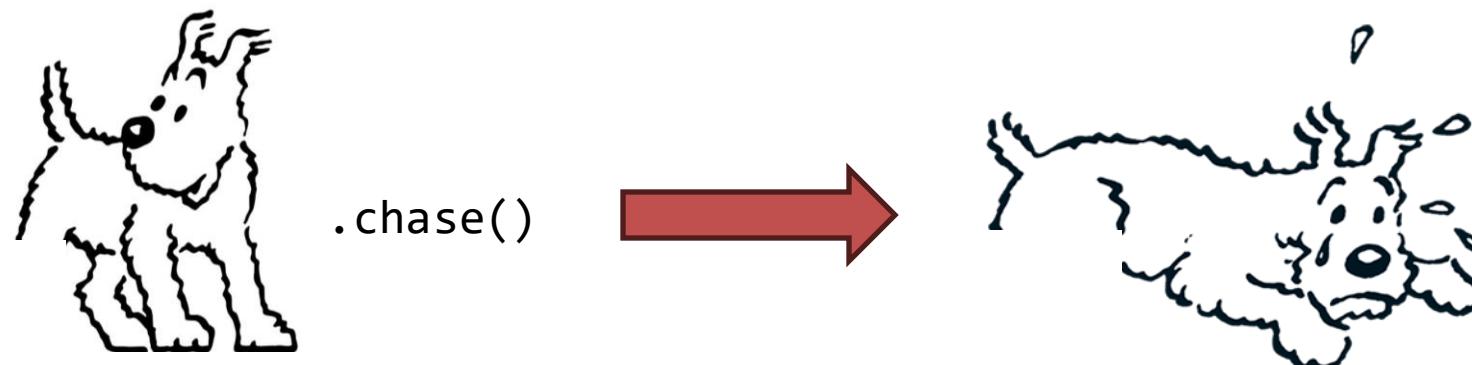


PP-vulnerable program

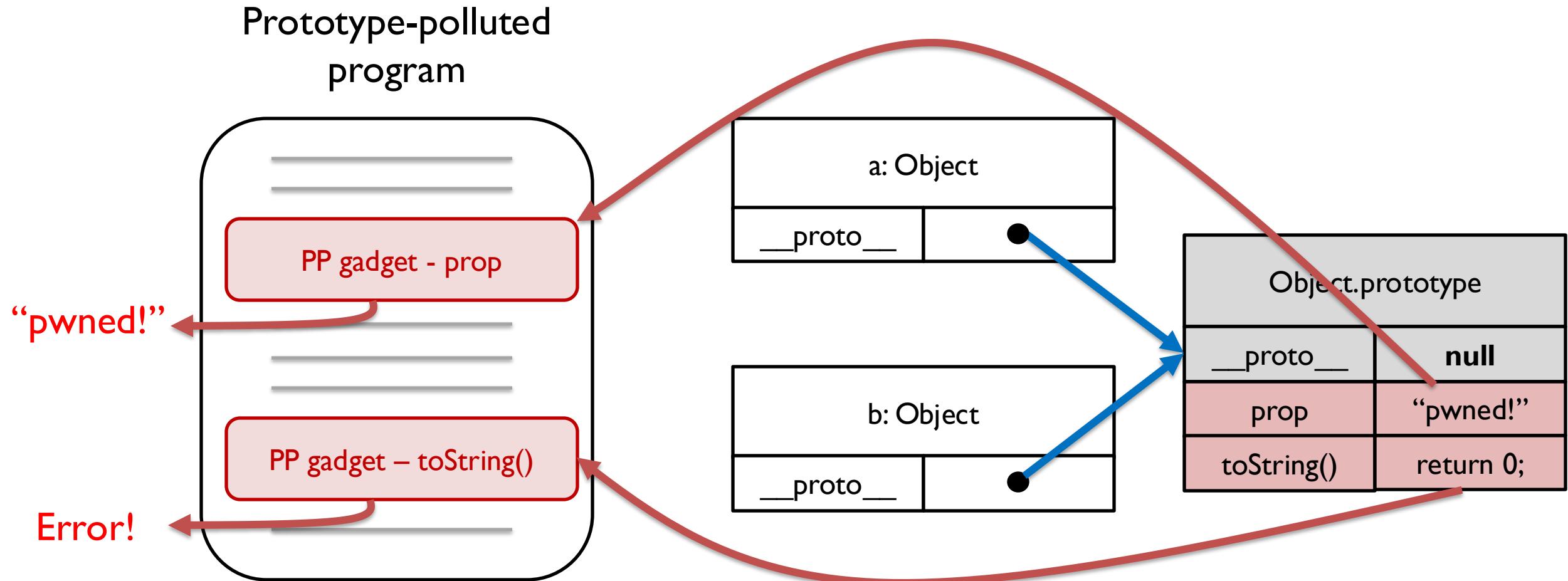


Gadget

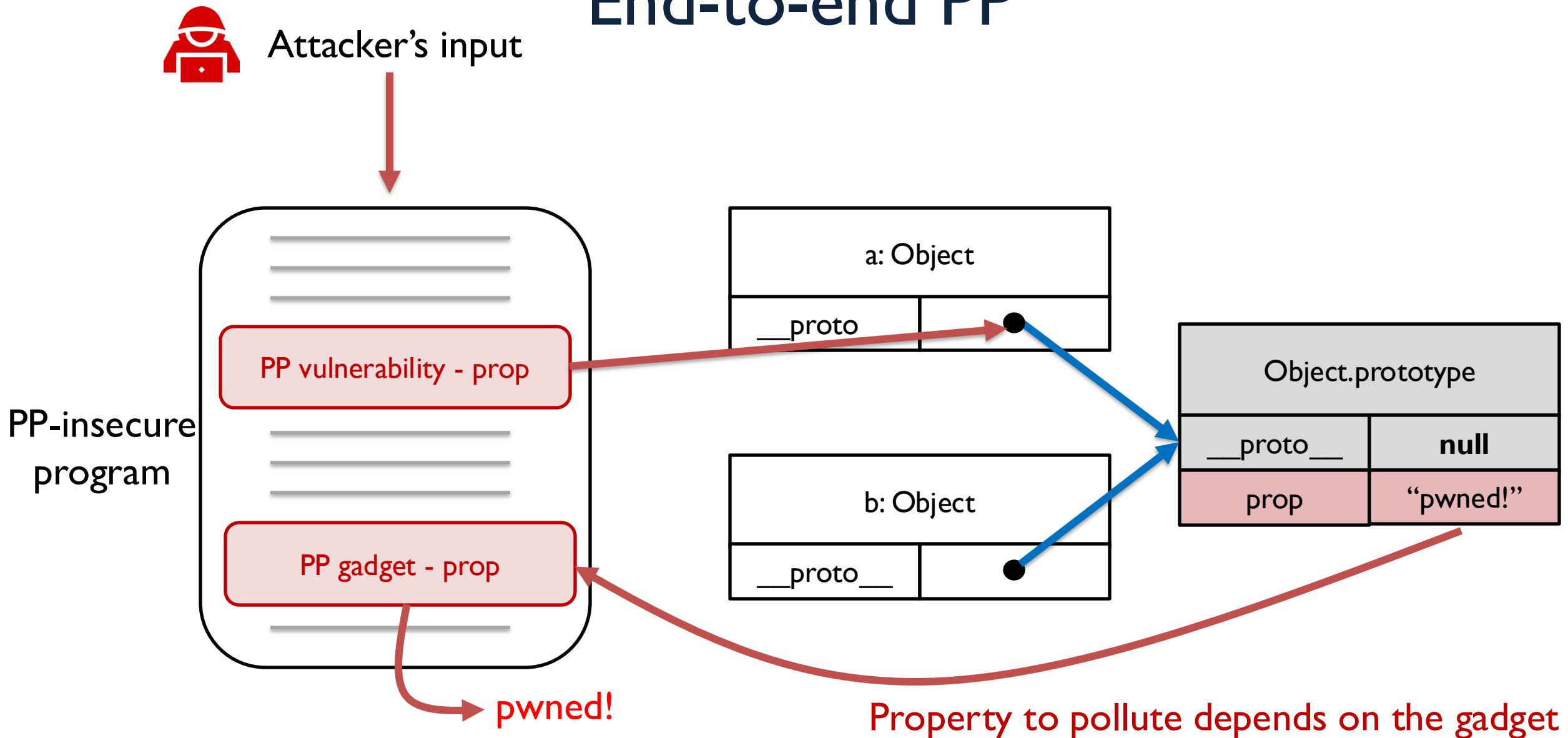
- PP's impact depends on the existence of **gadgets**
 - An **otherwise benign** piece of code which *inadvertently read from polluted properties* to **execute security-sensitive operations**
 - Triggering unintended behavior
- Examples: privilege elevation, reading secrets, RCE, DoS, log pollution



Gadget (cont.)



End-to-end PP



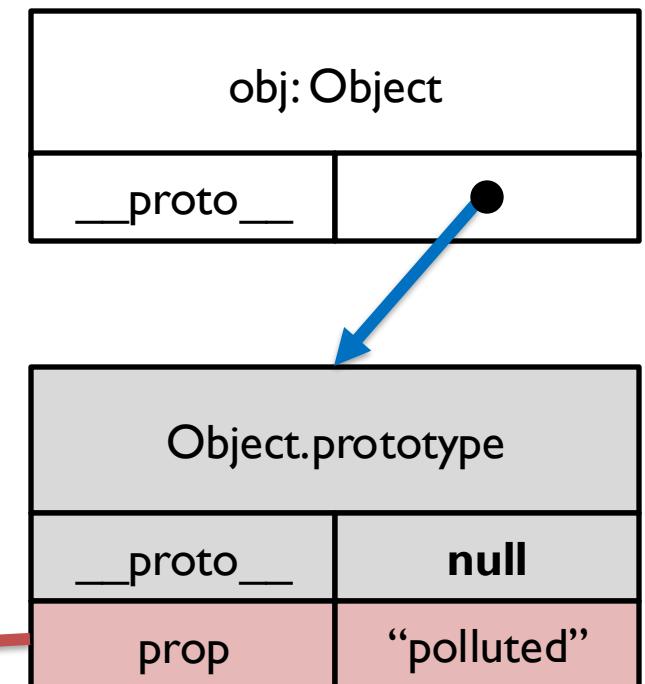
Example 1

```
var inProto = "__proto__"  
var inProp = "prop"  
var inVal = "polluted"  
  
var obj = {}  
var p = obj[inProto]  
p[inProp] = inVal  
  
console.log({}.prop) // "polluted"
```

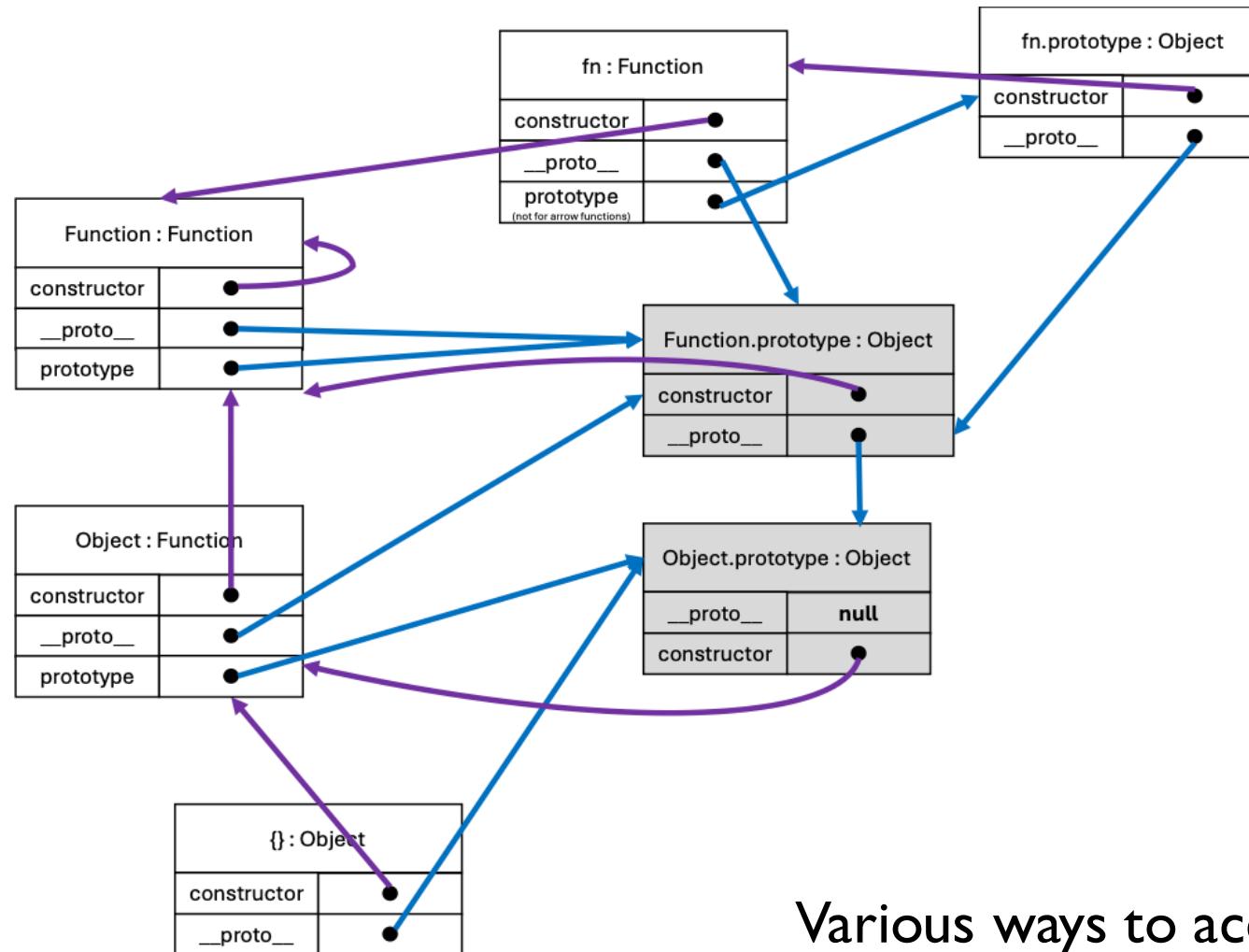
Malicious input

PP vulnerability

Gadget



Who says JavaScript is hard to learn?! 😊



Example 2

```
var inConstr = "constructor"
var inProto = "prototype"
var inProp = "prop"
var inVal = "polluted"

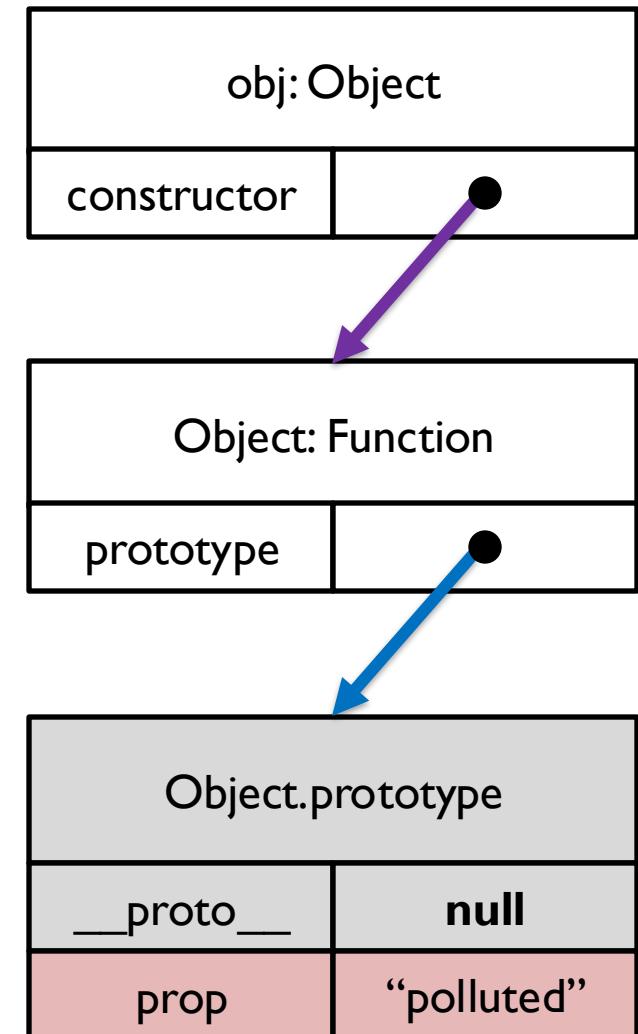
var obj = {}
var c = obj[inConstr]
var p = c[inProto]
p[inProp] = inVal

console.log({}.prop) // "polluted"
```

Malicious input

PP vulnerability

Gadget



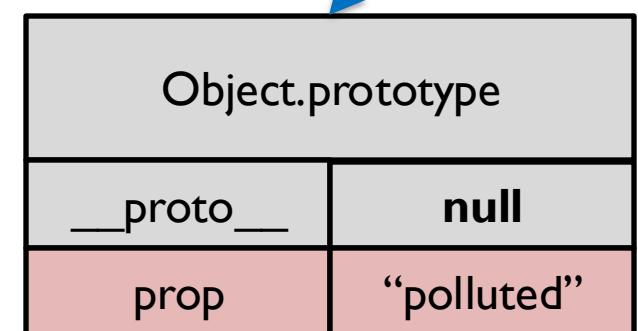
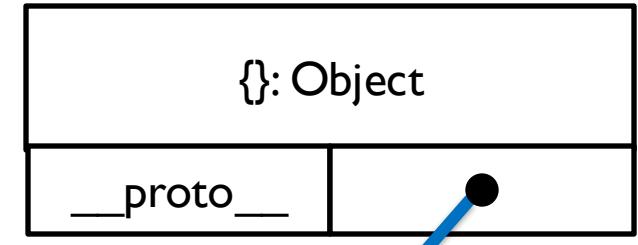
Example 3

```
function merge(dst, src) {  
  for (let key in src) {  
    if (!src.hasOwnProperty(key)) continue;  
    if (typeof dst[key] === 'object') {  
      merge(dst[key], src[key]);  
    } else {  
      dst[key] = src[key];  
    }  
  }  
  
  var input = '{"__proto__": {"prop": "polluted"}}';  
  merge({}, JSON.parse(input));  
  
  console.log({}.prop); // "polluted"
```

PP
vulnerability

Malicious
input
PP vulnerability

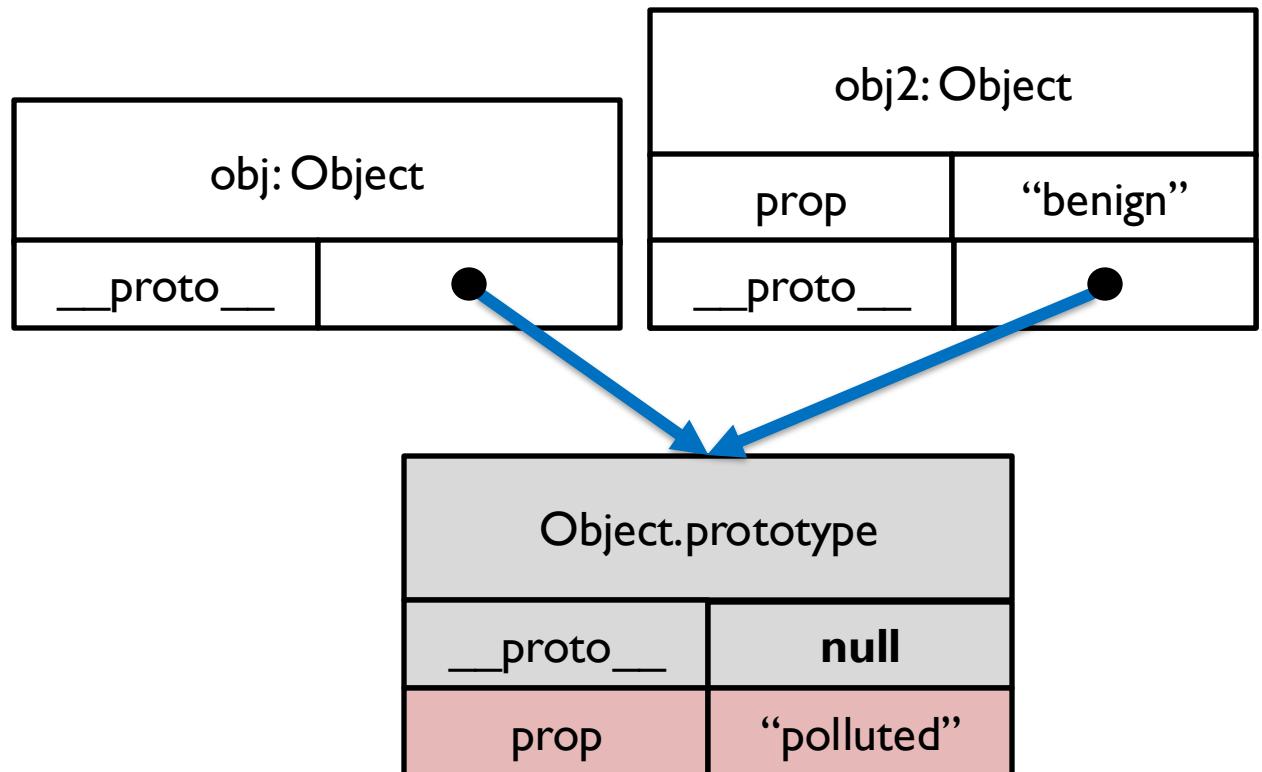
Gadget



Example 4

```
var inProto = "__proto__"  
var inProp = "prop"  
var inVal = "polluted"  
  
var obj = {}  
var p = obj[inProto]  
p[inProp] = inVal  
  
var obj2 = {}  
obj2.prop = "benign"  
  
console.log(obj2.prop) // "benign"
```

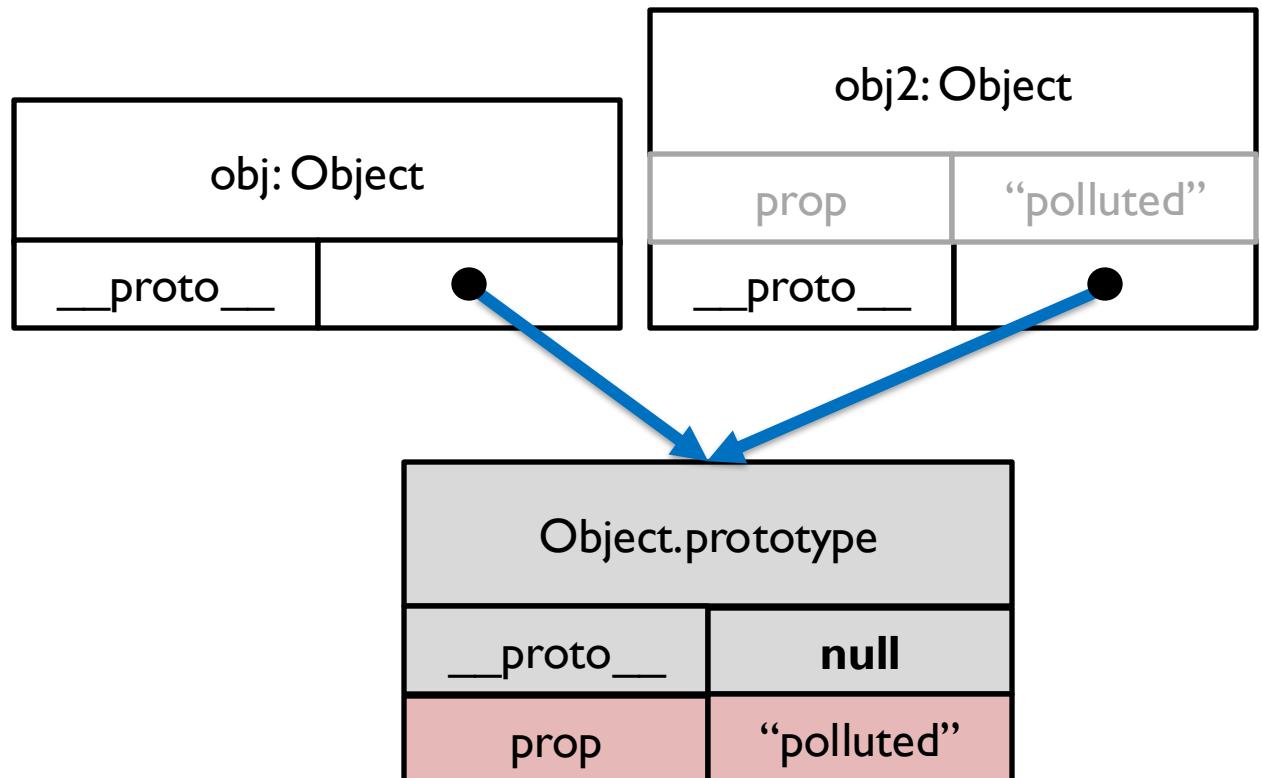
Observable behavior?



Example 5

```
var inProto = "__proto__"  
var inProp = "prop"  
var inVar = true  
  
var obj, obj2 = {}  
if (inVar)  
  obj[inProto][inProp] = "polluted"  
else  
  obj2.prop = "polluted"  
  
console.log(obj2.prop) // "polluted"
```

Observable behavior? 



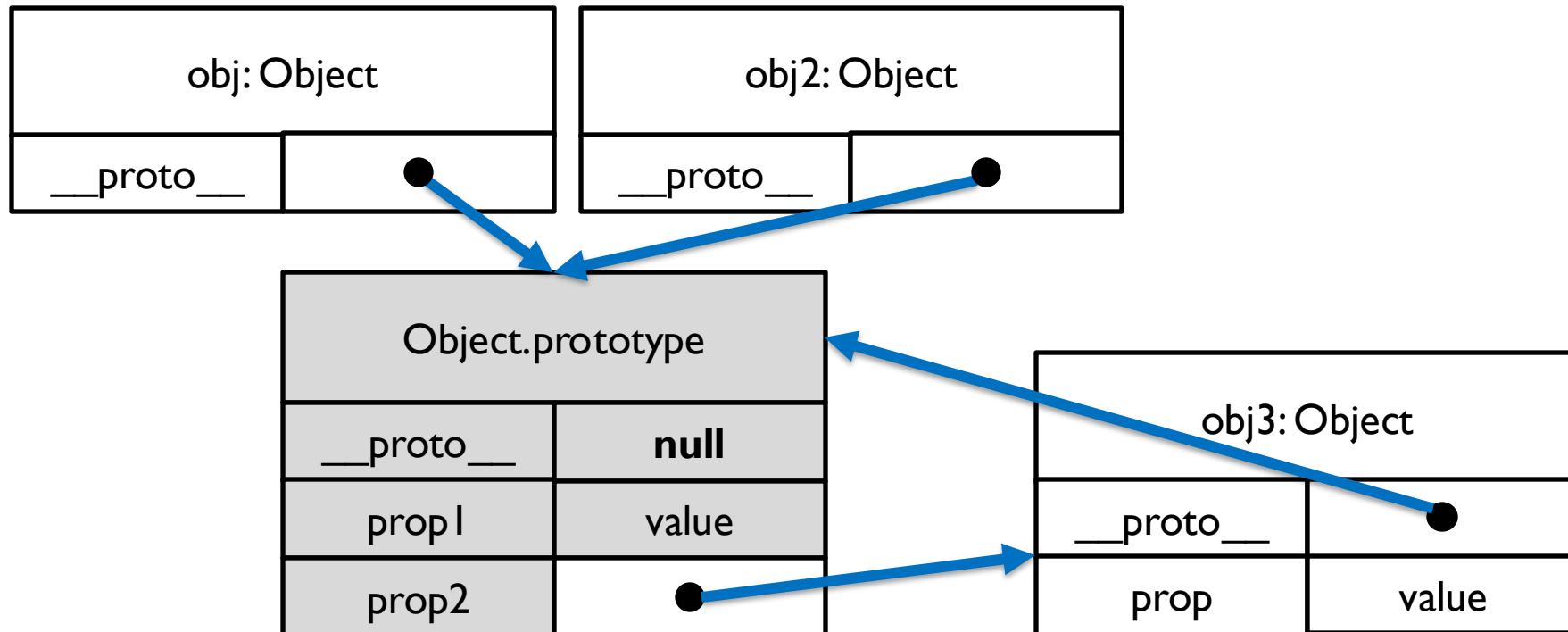
From PP to unintended behavior

- The literature focuses on *detection* rather than *defining* PP and gadgets *formally*
 - Any untrusted access to `__proto__` is forbidden (*shallow definition*)
 - Any modification to an object property that is reachable from a prototype object is forbidden (*deep definition*)
- Gadgets
 - ACE, SSRF, privilege escalation, DoS, log pollution, cryptographic downgrade, ...

What are the formal definitions of **PP-secure** and **gadget-free** programs?

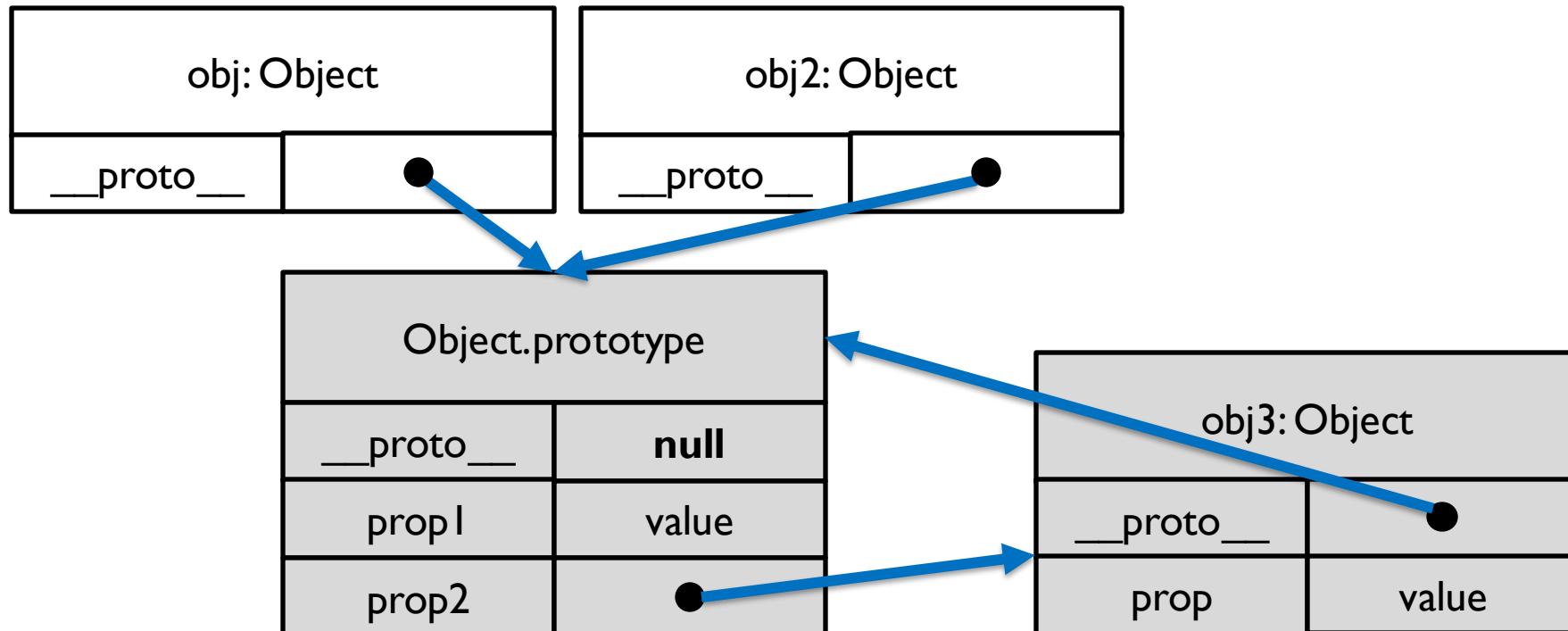
Shallow definition

The gray cells are expected to remain same for any untrusted input



Deep definition

The gray cells are expected to remain same for any untrusted input



Model language

A subset of the non-strict semantics of ECMA-262 standard

$v ::= s \mid n \mid b \mid \text{undefined} \mid \text{null}$

$e ::= v \mid x \mid e \oplus e \mid e(e) \mid \boxed{e[e] \mid \text{new } e(e)} \mid \text{function } (x) \ c$

$i ::= x \mid \boxed{i[e]}$

$c ::= \text{skip} \mid \boxed{i = e} \mid \text{if } (e) \ c \ \text{else } c \mid \text{while } (e) \ c \mid c; c \mid \text{return } e \mid \boxed{\text{out}_L(e)}$

PP security

PP-secure(c) $\stackrel{\text{def}}{=} \forall E_1, E_2, H_1, H_2.$

$(E_1, H_1) \simeq_T (E_2, H_2) \wedge$

$\Gamma \models \langle c, E_1, H_1, \emptyset \rangle \rightarrow^* \langle E_3, H_3, \tau_1 \rangle \wedge$

$\Gamma \models \langle c, E_2, H_2, \emptyset \rangle \rightarrow^* \langle E_4, H_4, \tau_2 \rangle$

$\implies \tau_1 \simeq_{stut}^{proto} \tau_2.$

Shallow vs.
deep

PP-Gadget freedom

$$\mathbf{PP\text{-}Gadget\text{-}Free}(c) \stackrel{\text{def}}{=} \forall E_1, E_2, H_1, H_2, H_3^P, H_4^P.$$

$$\Gamma \models \langle c, E_1, H_1, \emptyset \rangle \rightarrow_v^* \langle E'_1, H'_1, \tau_1 \rangle \wedge$$

$$\Gamma \models \langle c, E_2, H_2, \emptyset \rangle \rightarrow_{v'}^* \langle E'_2, H'_2, \tau_2 \rangle \wedge$$

$$v \sim_{\beta} v'$$

$$\implies$$

$$\Gamma \models \langle c, E_1, [H_1^O \uplus H_3^P], \emptyset \rangle \rightarrow_{v''}^* \langle E''_1, H''_1, \tau_3 \rangle \wedge$$

$$\Gamma \models \langle c, E_2, [H_2^O \uplus H_4^P], \emptyset \rangle \rightarrow_{v'''}^* \langle E''_2, H''_2, \tau_4 \rangle \wedge$$

$$v'' \sim_{\beta} v'''.$$

End-to-end PP

$E2E\text{-secure}(c) \stackrel{\text{def}}{=} \forall E_1, E_2, H_1, H_2.$

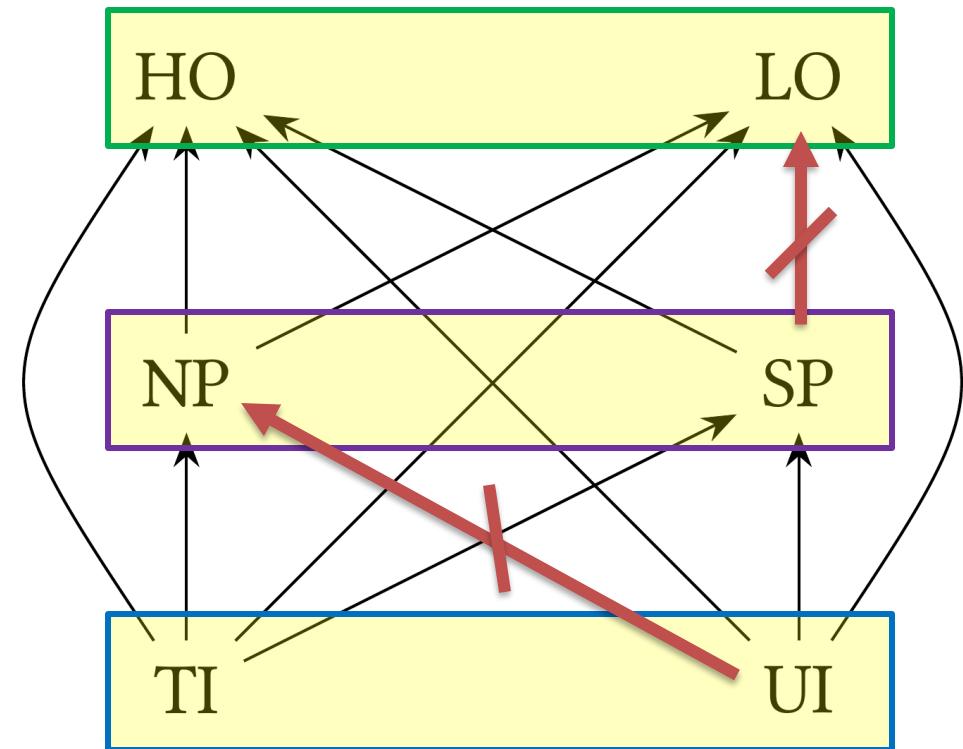
$(E_1, H_1) \simeq_T (E_2, H_2) \wedge$

$H_1 \simeq^{\kappa, NP} H_2 \wedge$

$\Gamma \models \langle c, E_1, H_1, \emptyset \rangle \rightarrow_v^* \langle E_3, H_3, \tau_1 \rangle \wedge$

$\Gamma \models \langle c, E_2, H_2, \emptyset \rangle \rightarrow_{v'}^* \langle E_4, H_4, \tau_2 \rangle$

$\implies v \sim_{\beta}^{LO} v' \vee \tau_1 \simeq_{stut}^{\kappa, SP} \tau_2.$



Ongoing work

- Sound enforcement mechanisms
 - Runtime monitoring
- Non-root pollution
- Degree of exploitability and attacker's knowledge
- Modeling mitigations
 - `Object.freeze()`, `Object.seal()`
 - Use of `new Set()`, `new Map()`



Takeaways

Formal modeling of prototype-pollution security

